

Notes on LCCC

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1. Introduction

1.1. Preliminary notions

1.1.1. Change of base functor

For this section, assume that we have a category \mathcal{C} which has pullbacks. Let X, Y be two objects in \mathcal{C} , and $f : X \rightarrow Y$. We can build

$$f^* : \mathcal{C}/Y \longrightarrow \mathcal{C}/X$$

the “base change” functor as follows: let $g : Z \rightarrow Y$ be an element of \mathcal{C}/Y .

$$\begin{array}{ccc} X \times_Y Z & \xrightarrow{\quad} & Z \\ \downarrow f^*(g) & \lrcorner & \downarrow g \\ X & \xrightarrow{\quad f \quad} & Y \end{array}$$

Furthermore, if $h : W \rightarrow Y$ is an other element of \mathcal{C}/Y and $\varphi : h \rightarrow g$, by pullback, there exists a unique $f^*(\varphi) : X \times_Y W \rightarrow X \times_Y Z$ making the following diagram commute

$$\begin{array}{ccccc} & & X \times_Y W & \xrightarrow{\quad} & W \\ & \swarrow f^*(\varphi) & \lrcorner & \searrow \varphi & \\ X \times_Y Z & \xrightarrow{\quad} & Z & & \\ \downarrow f^*(g) & \lrcorner & \downarrow f^*(h) & \searrow h & \\ X & \xrightarrow{\quad f \quad} & Y & & \end{array}$$

hence, $f^*(\varphi)$ is a morphism $f^*(h) \rightarrow f^*(g)$.

Lemma 1.1.1.1: f^* defines a functor.

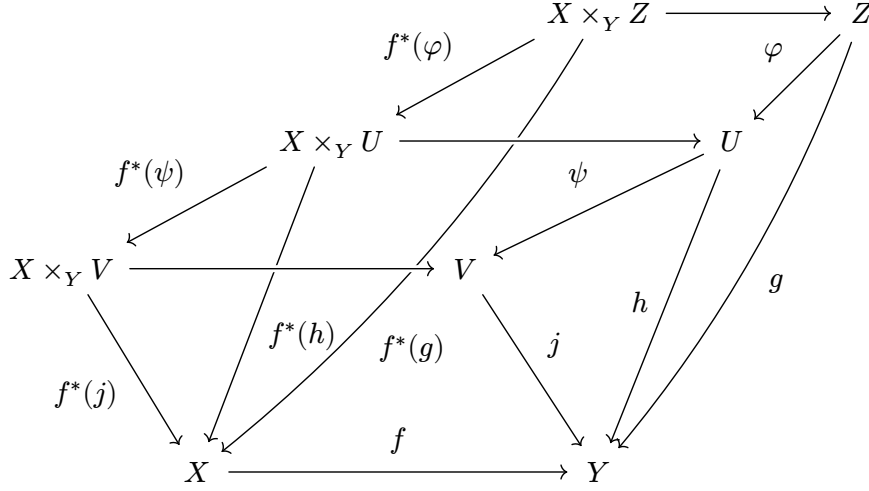
Proof: Let $(Z, g) : \mathcal{C}/Y$. Note that the following diagram commutes

$$\begin{array}{ccccc} & & X \times_Y Z & \xrightarrow{\quad} & Z \\ & \swarrow \text{id}_{f^*(g)} & \lrcorner & \searrow \text{id}_g & \\ X \times_Y Z & \xrightarrow{\quad} & Z & & \\ \downarrow f^*(g) & \lrcorner & \downarrow f^*(g) & \searrow g & \\ X & \xrightarrow{\quad f \quad} & Y & & \end{array}$$

thus $\text{id}_{f^*(g)}$ satisfies the universal property of $f^*(\text{id}_g)$, and so

$$f^*(\text{id}_g) = \text{id}_{f^*(g)}$$

Let $(U, h), (V, j) : \mathcal{C}/Y$, $\varphi : (Z, g) \rightarrow (U, h)$ and $\psi : (U, h) \rightarrow (V, j)$.



Note that the topmost outer rectangle commutes because both inner square commute, and that the leftmost, outermost triangle commutes too because both inner triangle commute, making $f^*(\psi) \circ f^*(\varphi)$ satisfy the same universal condition as $f^*(\psi \circ \varphi)$, so

$$f^*(\varphi \circ \psi) = f^*(\psi) \circ f^*(\varphi)$$

□

1.2. Main theorem

Definition 1.2.1 (Locally Cartesian Closed Category): A category \mathcal{C} is *locally cartesian closed* if, for any object $X : \mathcal{C}$, the category \mathcal{C}/X is cartesian closed.

Theorem 1.2.2: A category \mathcal{C} is locally cartesian closed if, and only if, it has

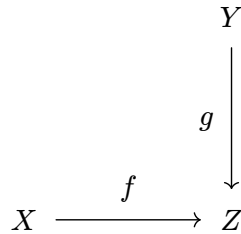
- pullbacks;
- for any morphism $f : X \rightarrow Y$ in \mathcal{C} , the functor f^* has a right adjoint Π_f , called the *dependent product at f* .

2. The direct part of the equivalence

Suppose we have a locally cartesian closed category \mathcal{C} .

2.1. Pullbacks

Let's show that \mathcal{C} has pullbacks. Let f, g be morphisms in \mathcal{C} spelling out the following diagram



Since \mathcal{C}/Z is cartesian closed, there exists a $(X \times_Z Y, h) : \mathcal{C}/Z$ the cartesian product of (X, f) and (Y, g) , with projections π_1 and π_2 :

$$\begin{array}{ccc}
X \times_Z Y & \xrightarrow{\pi_2} & Y \\
\pi_1 \downarrow & \searrow h & \downarrow g \\
X & \xrightarrow{f} & Z
\end{array}$$

This latter square (if we forget about h) is a pullback. Indeed, for any $U : \mathcal{C}$, $i : U \rightarrow X$ and $j : U \rightarrow Y$ making the following diagram commute

$$\begin{array}{ccc}
U & \xrightarrow{j} & Y \\
i \searrow & & \downarrow g \\
& X \times_Z Y & \xrightarrow{\pi_2} \\
& \pi_1 \downarrow & \\
& X & \xrightarrow{f} Z
\end{array}$$

Let $h' = g \circ j = f \circ i$. Note that $j : (U, h') \rightarrow (Y, g)$ and $i : (U, h') \rightarrow (X, f)$ in \mathcal{C}/Z , so by cartesianity, there exists a unique $\varphi : (U, h') \rightarrow (X \times_Z Y, h)$ such that $\pi_2 \circ \varphi = j$ and $\pi_1 \circ \varphi = i$, that is, the following diagram commutes

$$\begin{array}{ccc}
U & \xrightarrow{j} & Y \\
\varphi \searrow & & \downarrow g \\
& X \times_Z Y & \xrightarrow{\pi_2} \\
i \searrow & \pi_1 \downarrow & \\
& X & \xrightarrow{f} Z
\end{array}$$

Note that any φ making the two triangles commute is also a morphism $(U, h') \rightarrow (X \times_Z Y, h)$ in \mathcal{C}/Z , so φ is indeed unique in \mathcal{C} .

2.2. Dependent product

2.2.1. Definition

Let $f : X \rightarrow Y$ be a morphism in \mathcal{C} , let us define the functor

$$\Pi_f : \mathcal{C}/X \rightarrow \mathcal{C}/Y$$

Consider a $(Z, p) : \mathcal{C}/X$.

2.2.2. Adjunction